

# Intractability

CMPT 125

Mo Chen

SFU Computing Science

3/4/2020

# Lecture 34

Today:

- Finite Tile Puzzles
- Exponential-Time Algorithms
- NP-Complete Problems

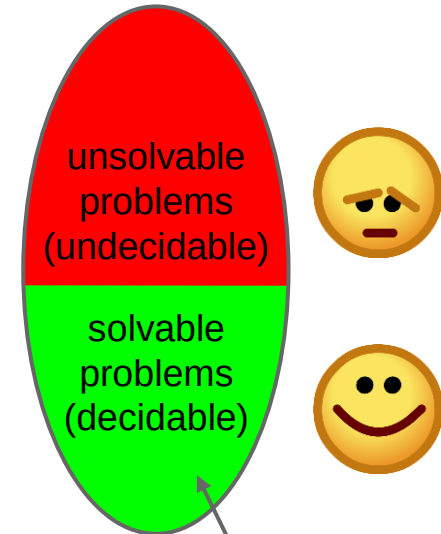
# Decidability (Review)

## Decision Problem:

- Answers a “Yes” / “No” question.

## Some problems don't have a solution

- called *undecidable*
- E.g., Tiling the plane
- E.g., Does program  $P$  have an infinite loop?
- E.g., Is program  $P$  correct?



## Some undecidable problems are “harder” than others

- Can write algorithms that:
  - are successful on restricted classes of inputs
  - work 99% of the time
- E.g., Lab grading server

These problems  
might not be that  
easy!

# Finite Tiling Puzzles

Problem: Given  $N = M^2$  tiles, can you tile an  $M \times M$  grid?

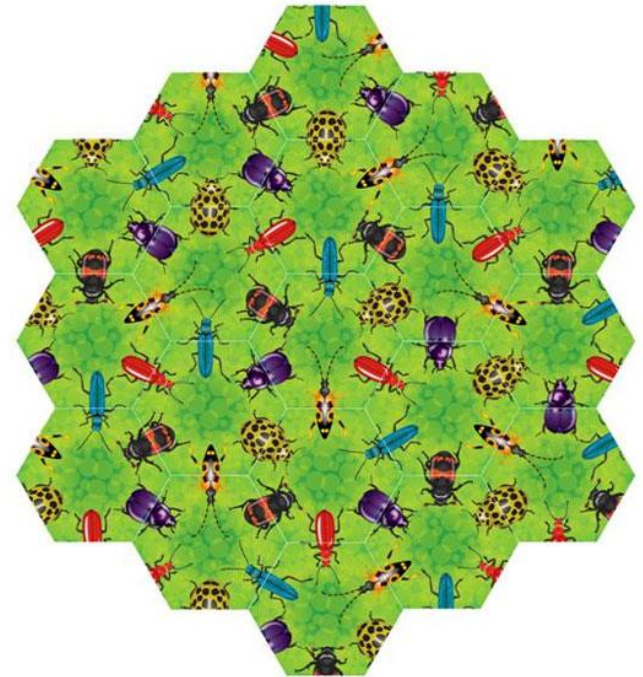
- harder version: allow to rotate / mirror

A brute force approach:

- Since there are a finitely many ways of arranging the tiles, and each arrangement can easily be tested for legality, try and test all arrangements
- decidable!

What's the running time?

- $N!$  arrangements means  $O(N!)$  time
- By the way,  $9! = 362880$



# Human Solution to Finite Tile Puzzle

You would also use brute-force, but add one tile at a time.

- If tile doesn't fit, then try another.
- If all tries lead to failure, then remove the previous tile.

Algorithm is called *backtracking*.

- recursive
- generates partial solutions
- we didn't do  $8!$  steps, because we rejected many permutations early

What makes a “hard” puzzle?

- many partial solutions, but . . .
- few correct solutions (usually one)



# Exponential vs Polynomial Time

$N!$  is the fastest growing function yet

- faster than any polynomial

Fastest known algorithm to solve every bug puzzle is  $2^N$ , which also grows very fast

- but remember: many puzzles can be solved quickly in practice!

**Rule of Thumb: A typical computer will do 1 billion operations in around 1 second.**

Q. What's the maximum practical  $N$ ?

|          | $N$                  | $N \log N$           | $N^2$             | $N^3$             | $2^N$ | $N!$ |
|----------|----------------------|----------------------|-------------------|-------------------|-------|------|
| 1 second | $10^9$               | $4.0 \times 10^7$    | $3.1 \times 10^4$ | $10^3$            | 30    | 12   |
| 10 years | $3.2 \times 10^{17}$ | $6.0 \times 10^{15}$ | $5.6 \times 10^8$ | $6.8 \times 10^5$ | 58    | 19   |

# Reasonable vs Unreasonable Time

The functions  $2^N$  and  $N!$  easily dwarf the growth of all functions of the form  $N^k$ , for any fixed  $k$ .

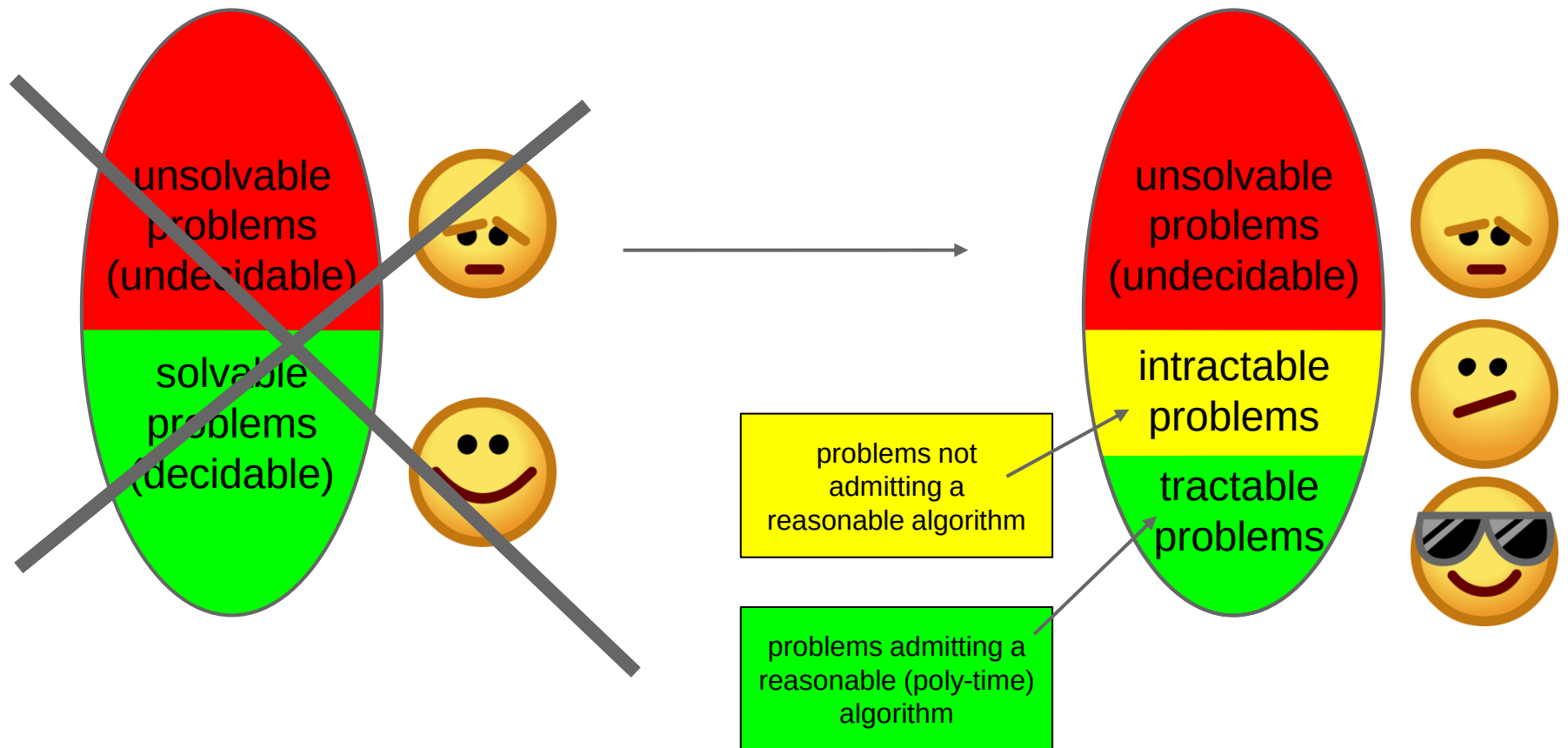
Two provisos:

- $N^{1000}$  also grows really stinking fast
- There are some linear algorithms with massive leading constants

But for the most part, the distinction is valid:

- “good” is polynomial — *tractable*
- “bad” is super-polynomial — *intractable*

# The World of All Problems





# Intractability and NPC

Q. Is the finite tile puzzle intractable or tractable?

Mathematically speaking, it is in no man's land.

- no one has proven that exponential time is required
  - exponential lower bounds have been proved for some problems (but not this one)
- no polynomial-time algorithm has yet been found
  - backtracking is one algorithm that works well on most bug puzzles, but not all of them.

The finite tile puzzle belongs to a class called *NPC* — the *NP*-Complete problems.

# Other Problems in *NPC*

There are many natural problems which are similar to the finite tile puzzle.

They come from a variety of domains.

- **The Travelling Salesman Problem (TSP)**

Given a road network connecting  $N$  cities, plan the fastest route that passes through all  $N$  of them.

- **Subset-sum.**

Given a list of  $N$  numbers and a target number  $t$ , find a subset of those numbers that sums to  $t$ .

# More Examples

- **Knapsacking**

Given a list of  $N$  items of weight  $w_1, w_2, \dots, w_N$  and value  $v_1, v_2, \dots, v_N$ , pack a car whose maximum load is  $W$  such that value is maximized

- **Scheduling**

Given a list of  $N$  students each of which has up to 5 final exams, devise the minimum schedule so that no exams overlap for any students

- **Satisfiability**

Given a logical expression of length  $N$ , find a true/false substitution that will yield “true”

# NPC — Rising and Falling Together

There are several hundred problems sharing remarkable properties:

- best known algorithm is exponential
- best lower bound is polynomial
- if one is intractable, then they all are, but . . .
- if one is tractable, then they all are!

$P \stackrel{?}{=} NP$  (Cook-Levin 1971)

- The most important open problem in CS
- Also considered a major open problem in Mathematics
- <https://youtu.be/YX40hbAHx3s>

# A Use for Intractability

Sometimes the bad news can be used constructively:

- in cryptography and security

**General Strategy:** Devise a cryptosystem so that unauthorized decryption is expensive.

Most public-key crypto relies on large primes

- you can eavesdrop only if you can factor extremely large numbers efficiently
- integer factorization is believed to be intractable

# CMPT 125 — Topics Covered

## Algorithms:

- measuring performance
- the worst case  $\rightarrow$  big- $O$
- software engineering principles
- brute-force paradigm
- sorting and searching
- assertions, pre/post-conditions, invariants, invariant proofs
- divide & conquer paradigm
- recursion & recursive invariants
- ADTs: stacks, queues
- linked lists, rooted trees
- binary search trees
- regular expressions & FSMs
- floating point encoding
- undecidability, intractability

## Coding:

- declare all variables
- pass-by-value
- arrays are fixed length
- strings
- pointers
- coding style
- recursion
- struct
- interfaces
- ADTs
- linked data structures
- struct  $\rightarrow$  class
- templates & the STL