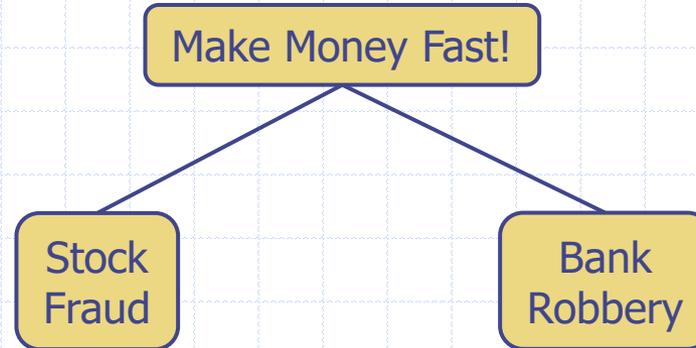


Binary Trees

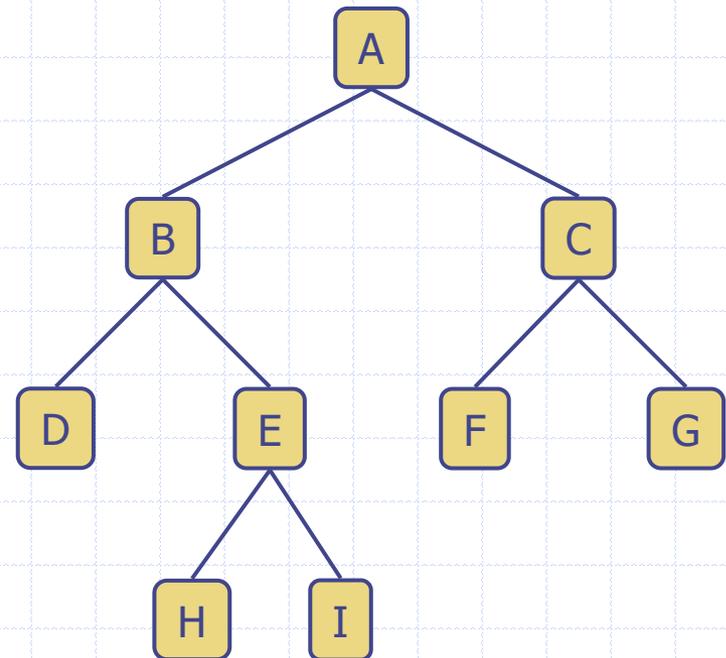
Section 7.3



Binary Trees

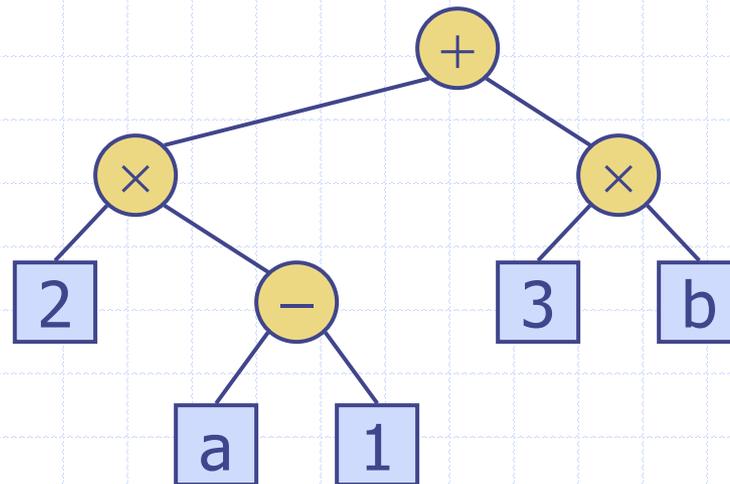
- A binary tree is a tree with the following properties:
 - Each internal node has at most two children (exactly two for **proper** binary trees)
 - The children of a node are an ordered pair
- We call the children of an internal node **left child** and **right child**
- Alternative recursive definition: a binary tree is either
 - a tree consisting of a single node, or
 - a tree whose root has an ordered pair of children, each of which is a binary tree

- Applications:
 - arithmetic expressions
 - decision processes
 - searching



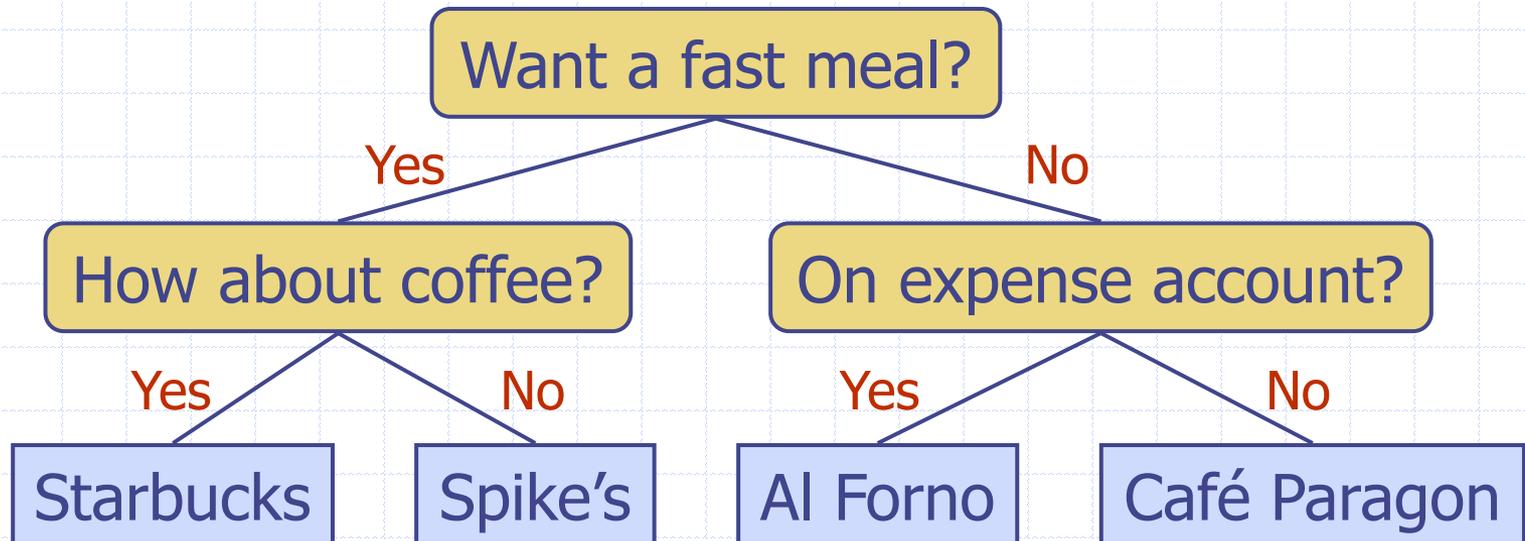
Arithmetic Expression Tree

- Binary tree associated with an arithmetic expression
 - internal nodes: operators
 - external nodes: operands
- Example: arithmetic expression tree for the expression $(2 \times (a - 1) + (3 \times b))$



Decision Tree

- Binary tree associated with a decision process
 - internal nodes: questions with yes/no answer
 - external nodes: decisions
- Example: dining decision



Properties of Binary Trees

□ Notation

n number of nodes

e number of external nodes

i number of internal nodes

h height

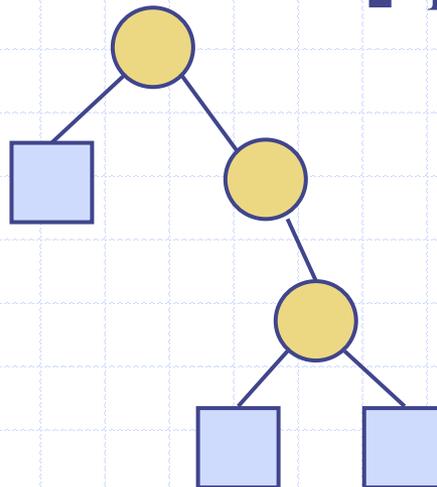
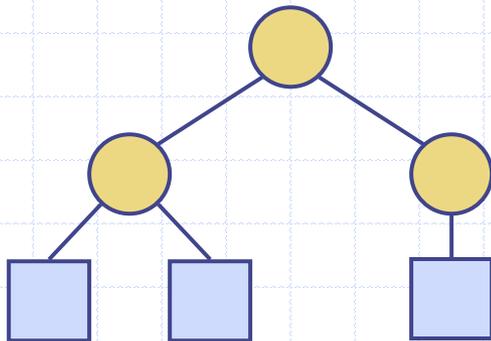
◆ Properties:

- $h + 1 \leq n \leq 2^{h+1} - 1$

- $1 \leq e \leq 2^h$

- $h \leq i \leq 2^h - 1$

- $\log(n + 1) - 1 \leq h \leq n - 1$



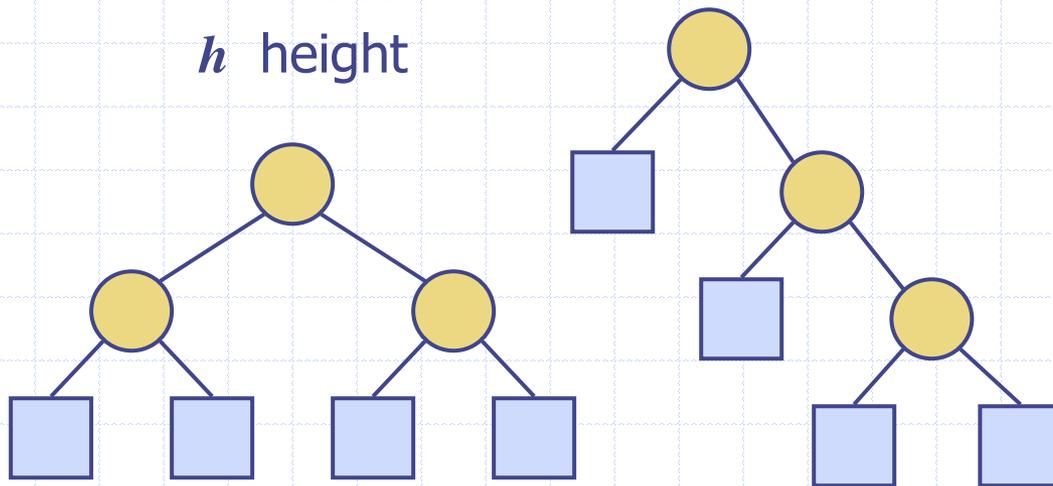
Properties of Proper Binary Trees

□ Notation

- n number of nodes
- e number of external nodes
- i number of internal nodes
- h height

◆ Properties:

- $2h + 1 \leq n \leq 2^{h+1} - 1$
- $h + 1 \leq e \leq 2^h$
- $\log(n + 1) - 1 \leq h \leq (n - 1)/2$



- $e = i + 1$
- $n = 2e - 1$
- $h \leq i$
- $h \geq \log_2 e$

BinaryTree ADT

- The BinaryTree ADT extends the Tree ADT, i.e., it inherits all the methods of the Tree ADT
- Additional methods:
 - position `p.left()`
 - position `p.right()`
- Update methods may be defined by data structures implementing the BinaryTree ADT

BinaryTree ADT

- There is a **position** class associated with the tree, to provide public access to nodes.
- Tree::Position supports member functions
 - **element()** Return the associated element.
 - **left()** Return the left child.
Error if this is an external node.
 - **right()** Return the right child.
Error if this is an external node.
 - **parent()** Return the parent. Error if this is the root.
 - **isRoot()** Return **true** if this is the root; else **false**.
 - **isExternal()** Return **true** if this is external; else **false**.

BinaryTree ADT

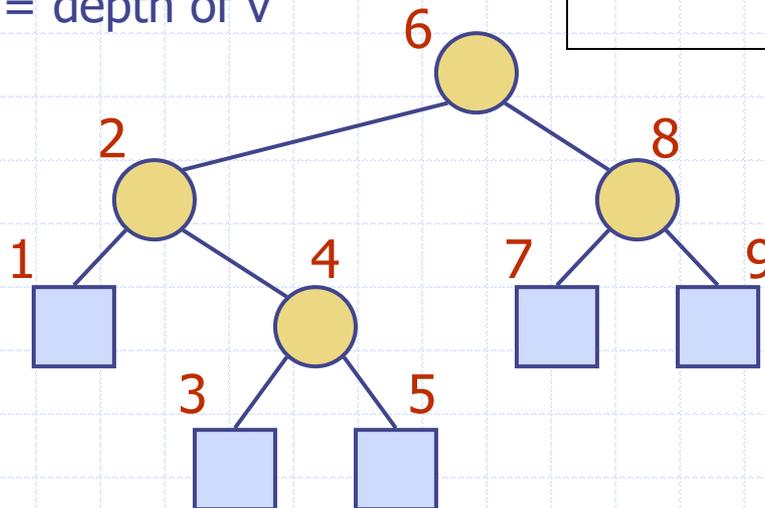
- The Tree itself supports member functions:
 - `size()` Return the number of nodes in the tree.
 - `empty()` Return **true** if the tree is empty; else **false**.
 - `root()` Return a position for the tree's root.
Error if the tree is empty.
 - `positions()` Return a position list of all of the nodes of the tree.

- Note that we haven't defined update functions, so building a tree isn't possible right now. We'll do that later.

Inorder Traversal

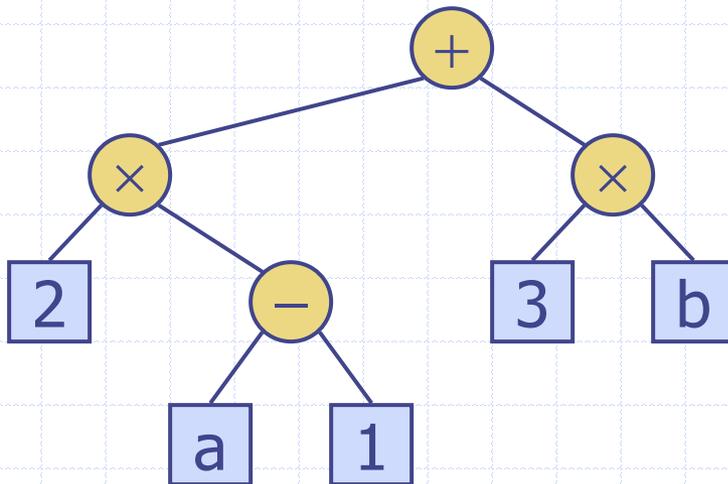
- In an inorder traversal a node is visited after its left subtree and before its right subtree
- Application: draw a binary tree
 - $x(v)$ = inorder rank of v
 - $y(v)$ = depth of v

```
Algorithm inOrder( $v$ )  
  if  $\neg v.isExternal()$   
    inOrder( $v.left()$ )  
  visit( $v$ )  
  if  $\neg v.isExternal()$   
    inOrder( $v.right()$ )
```



Print Arithmetic Expressions

- Specialization of an inorder traversal
 - print operand or operator when visiting node
 - print "(" before traversing left subtree
 - print ")" after traversing right subtree



Algorithm *printExpression(v)*

if $\neg v.isExternal()$

print("(")

printExpression(*v.left*())

print(*v.element*())

if $\neg v.isExternal()$

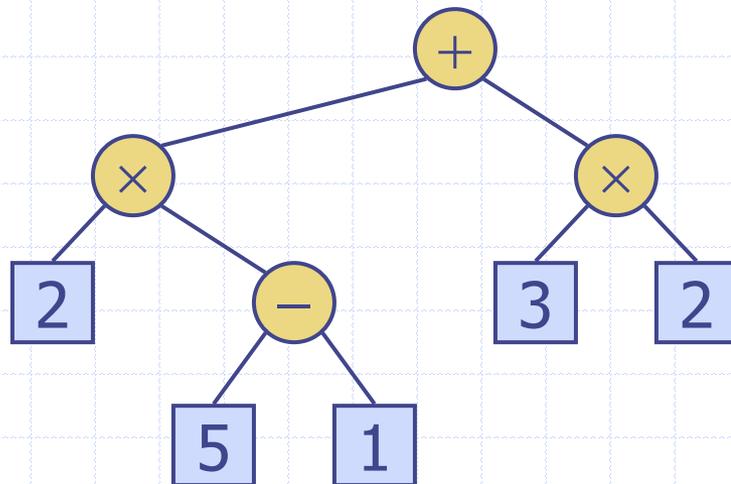
printExpression(*v.right*())

print(")")

$((2 \times (a - 1)) + (3 \times b))$

Evaluate Arithmetic Expressions

- Specialization of a **postorder** traversal
 - recursive method returning the value of a subtree
 - when visiting an internal node, combine the values of the subtrees



Algorithm *evalExpr(v)*

if *v.isExternal()*

return *v.element()*

else

x ← *evalExpr(v.left())*

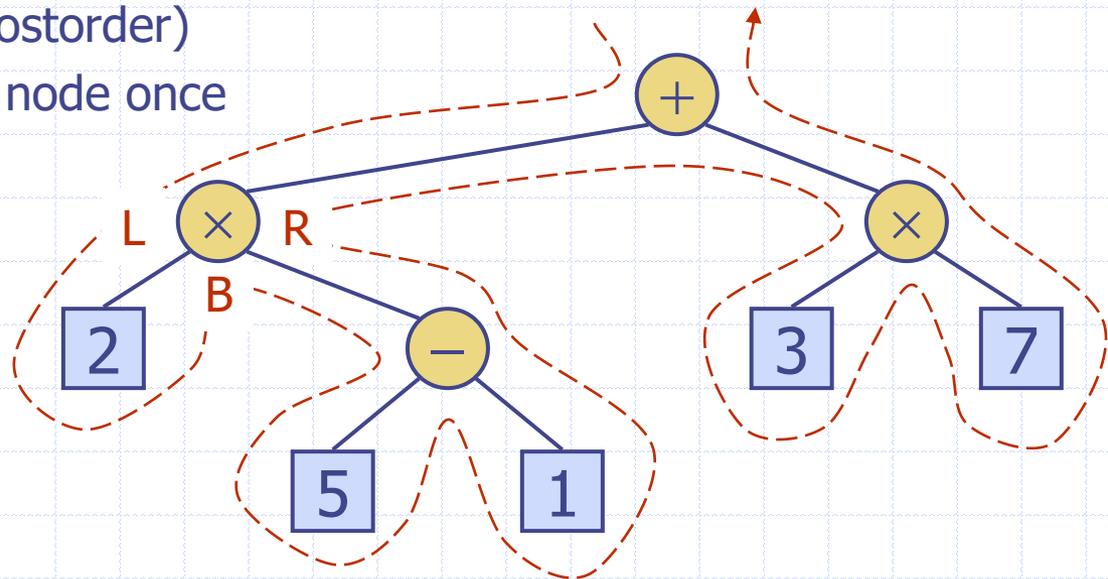
y ← *evalExpr(v.right())*

◇ ← operator stored at *v*

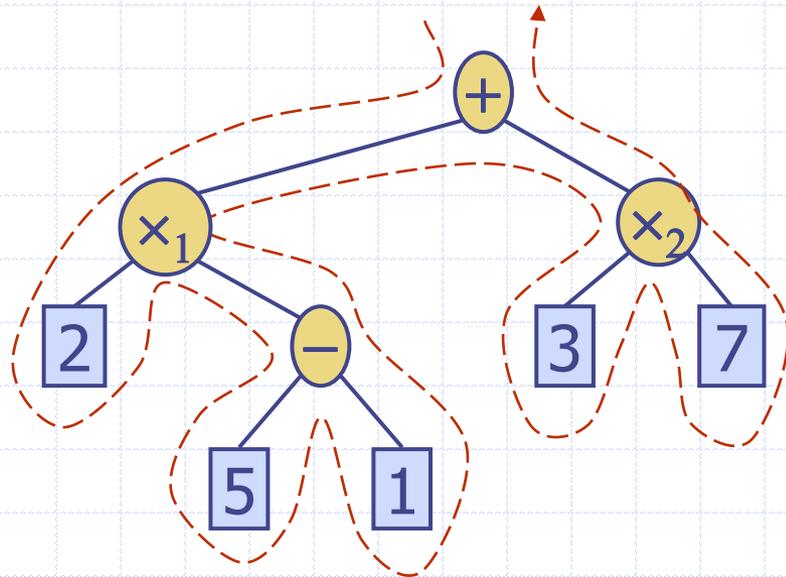
return *x* ◇ *y*

Euler Tour Traversal

- Generic traversal of a binary tree
- Includes a special cases the preorder, postorder and inorder traversals
- Walk around the tree and visit each internal node three times:
 - on the left (preorder)
 - from below (inorder)
 - on the right (postorder)
- Visit each external node once



Euler Tour Traversal

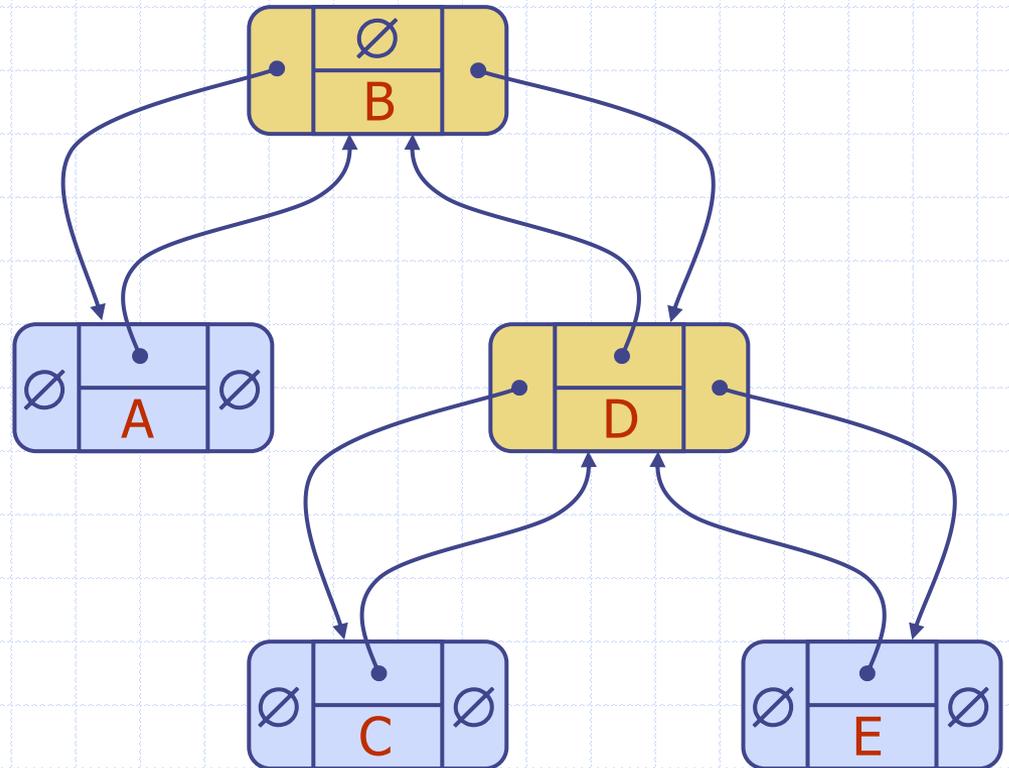
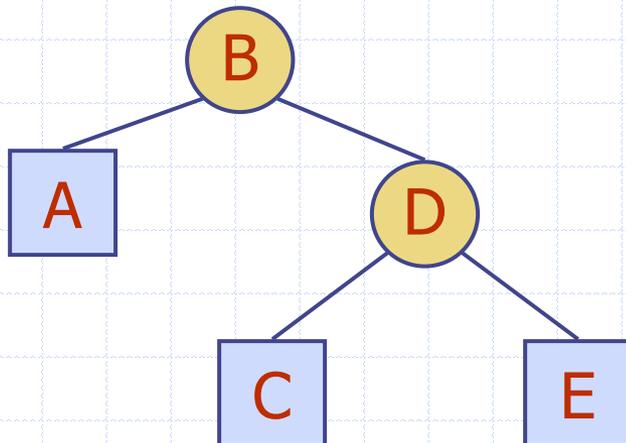


visitLeft(+)
visitLeft(x₁)
visit(2)
visitBottom(x₁)
visitLeft(-)
visit(5)
visitBottom(-)
visit(1)
visitRight(-)
visitRight(x₁)
visitBottom(+)
visitLeft(x₂)
visit(3)
visitBottom(x₂)

visit(7)
visitRight(x₂)
visitRight(+)

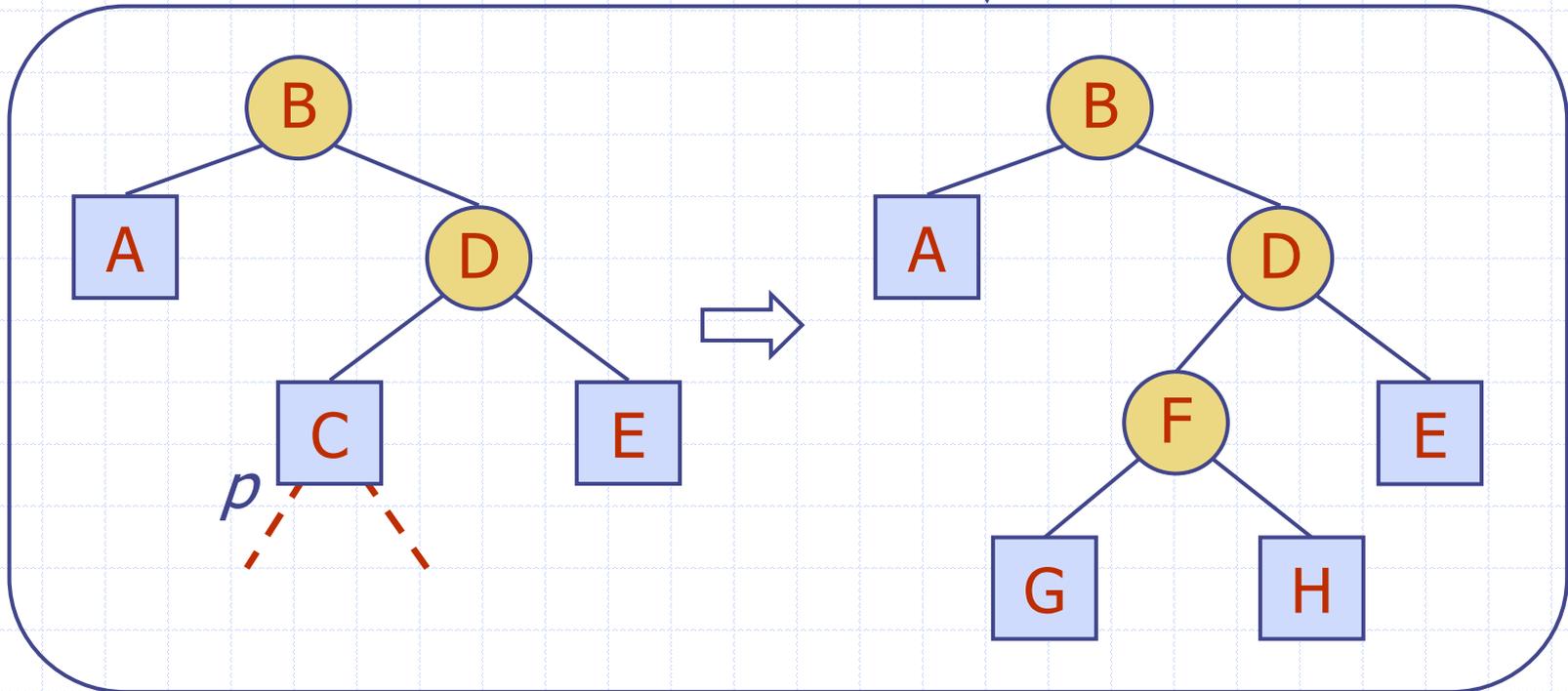
Linked Structure for Binary Trees

- A node is represented by an object storing
 - Element
 - Parent node
 - Left child node
 - Right child node
- Node objects implement the Position ADT



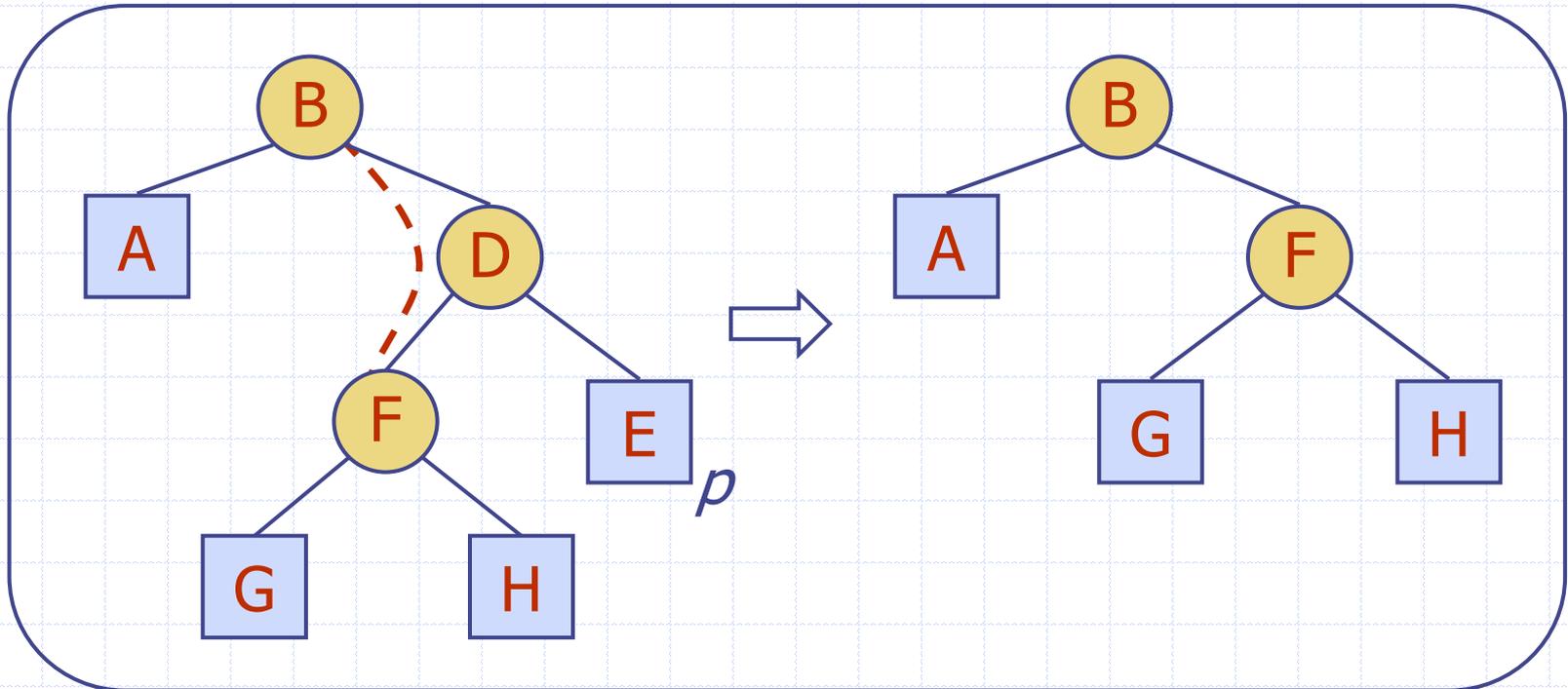
Proper Binary Tree Updates

- ❑ createRoot()
- ❑ expandExternal(p)
- ❑ removeAboveExternal(p)



Proper Binary Tree Updates

- removeAboveExternal(p)



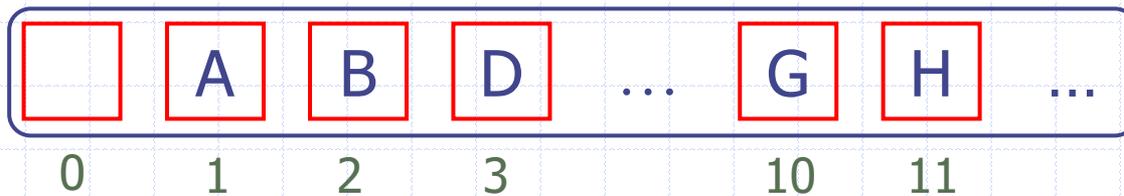
Binary Tree Performance Using Linked Structure

<i>Operation</i>	<i>Time</i>
left, right, parent, isExternal, isRoot	$O(1)$
size, empty	$O(1)$
root	$O(1)$
expandExternal, removeAboveExternal	$O(1)$
positions	$O(n)$

- Space usage is $O(n)$.

Array-Based Representation of Binary Trees

- Nodes are stored in an array A



- Node v is stored at $A[\text{rank}(v)]$
 - $\text{rank}(\text{root}) = 1$
 - if node is the left child of $\text{parent}(\text{node})$,
 $\text{rank}(\text{node}) = 2 \cdot \text{rank}(\text{parent}(\text{node}))$
 - if node is the right child of $\text{parent}(\text{node})$,
 $\text{rank}(\text{node}) = 2 \cdot \text{rank}(\text{parent}(\text{node})) + 1$

