Binary Trees

CMPT 125
Mar. 18

Lecture 27

Today:

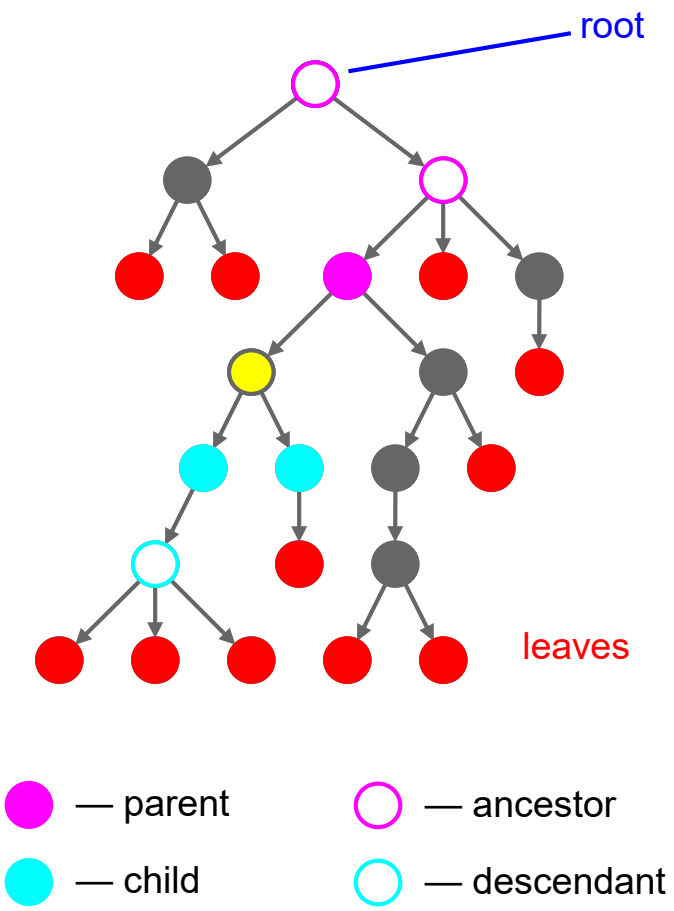
- Binary Trees
- Recursive Definitions of Trees
- Binary Tree Implementation
- Expression Trees
- Traversals
- Grammars

Rooted Trees (Review)

A *rooted tree* is a tree where all but one vertex has exactly one inbound edge (from its *parent*).

- usually drawn by level, top down
- *root* vertex has no inbound edge
- *leaf* vertex has no outbound edge
- parents point to *children*
- *ancestors* point to *descendants* via a downward path

A *binary tree* is a rooted tree in which no vertex has more than 2 children.



Subtrees and Recursive Definitions

There are many *subtrees* within a binary tree:

- the two most important are the left and right subtrees
- rooted at the left and right children of the root
- to visualize, remove the root!

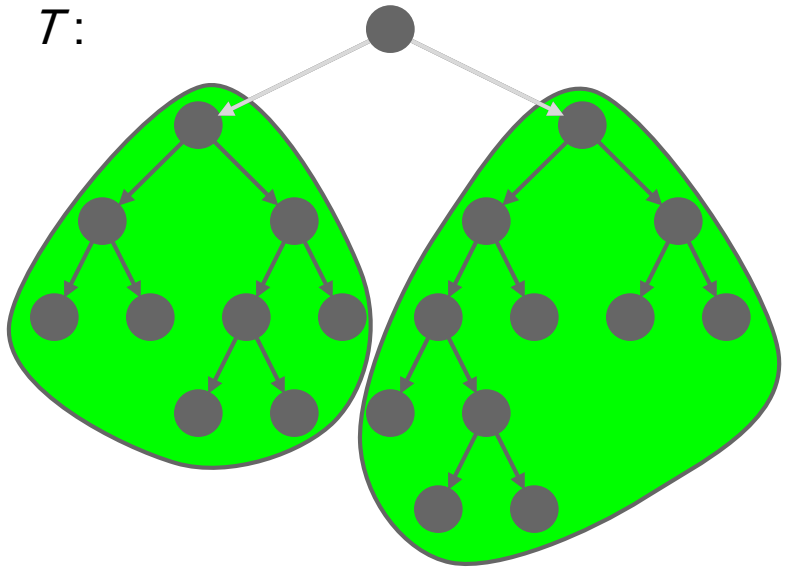
Leads to a recursive definition:

T is a binary tree when:

- *T* is an empty tree (i.e., no vertices)

  OR . . .

- *T* has a root vertex whose left and right subtrees are binary trees



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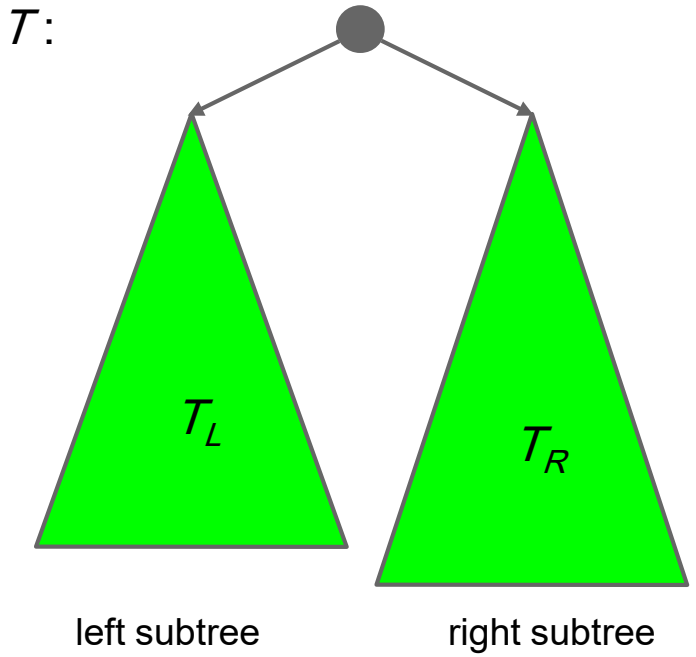


Trees and Recursion

Recursive definitions benefit you in two ways:

- allow you to write recursive algorithms
- allow you to reason about the structure by recursion (or induction)

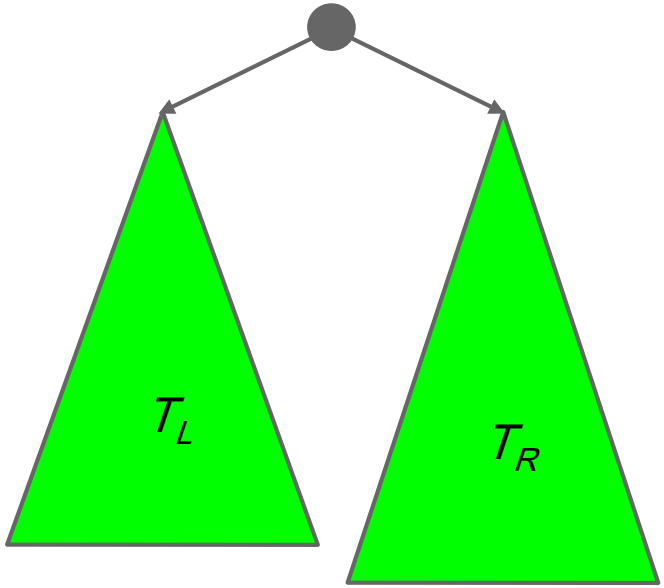
How to build a binary tree in C++?

- use the recursive definition
- adopt similar strategy to a linked list

```
struct LLnode {
    int data;
    struct LLnode * next;
};
```

T is a binary tree when either:

- *T* is an empty tree
- *T* has a root vertex whose left and right subtrees are binary trees



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struct BTnode {
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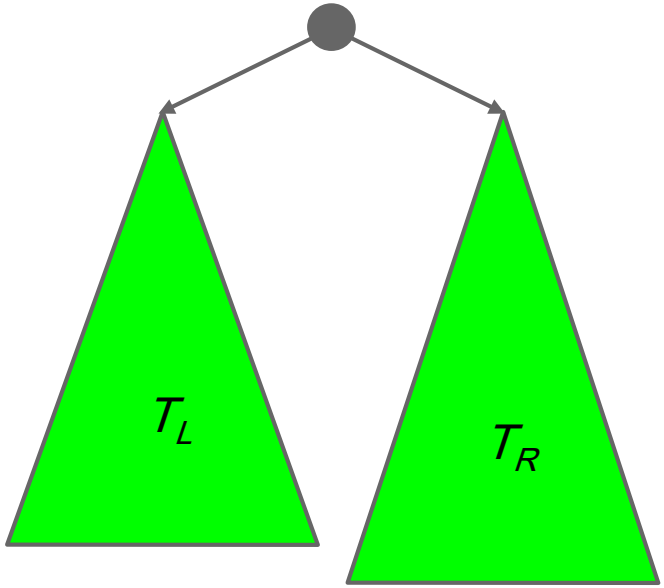

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struct BTnode {
    int data;
    struct BTnode * left;
    struct BTnode * right;
};
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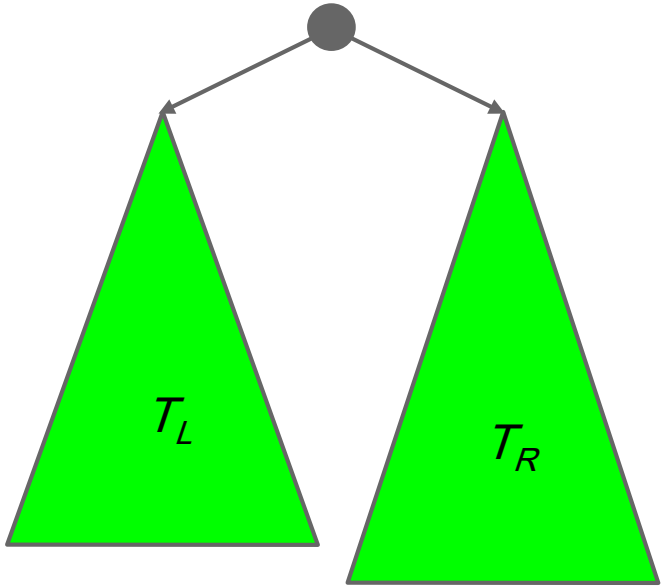

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struct BTnode {
    int data;
    struct BTnode * left;
    struct BTnode * right;
    struct BTnode * parent;
};
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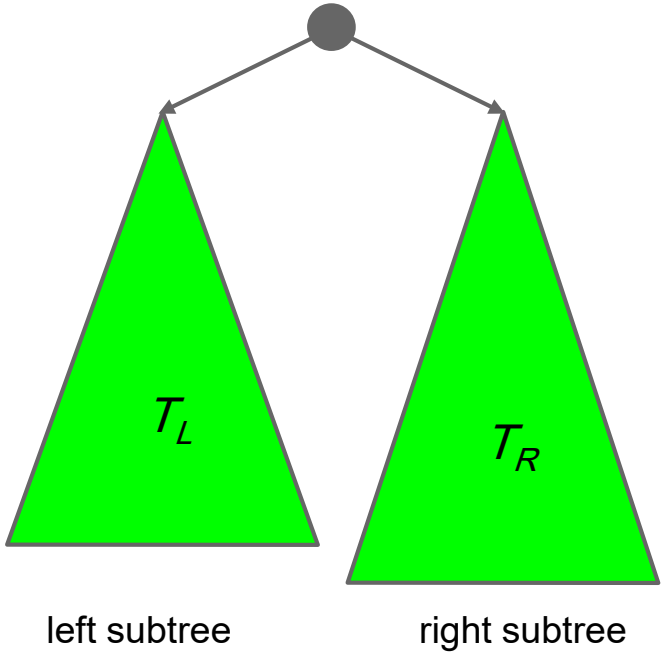


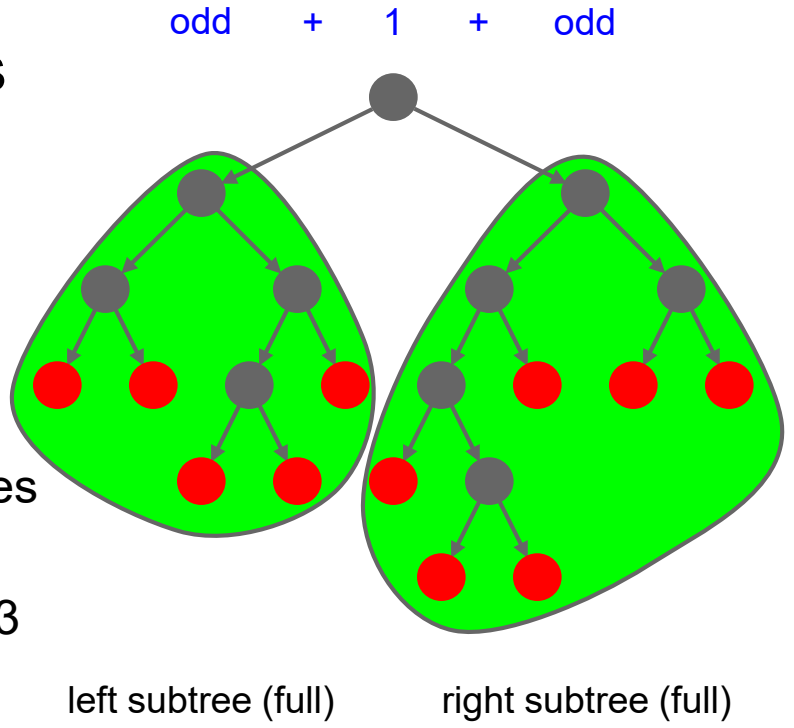
Reasoning about Trees

A *full binary tree* is a non-empty binary tree, where each vertex has exactly 0 or 2 children.

Theorem: A full binary tree always has an odd number of vertices.

Proof by induction:

- If the root has 0 children, then the tree has only one vertex, which is odd.
- If the root has 2 children, then their subtrees must also be full, and by induction, odd. The total number of vertices is the sum of 3 odd numbers, which is odd.

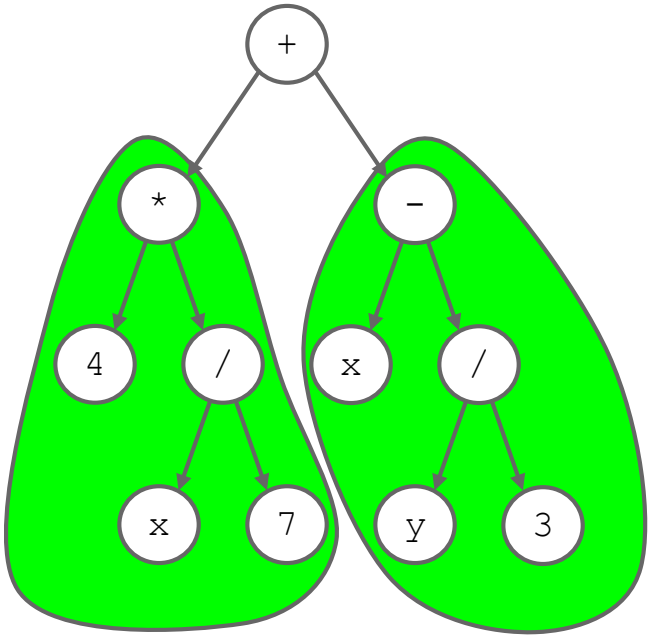


Expression trees are full.

Expression Trees

An *expression tree* is a full binary tree that represents an arithmetic calculation:

- internal nodes are binary operators
- leaves are numbers



left exp'n tree + right exp'n tree

Expression Trees and Postfix

An *expression tree* is a full binary tree that represents an arithmetic calculation:

- internal nodes are binary operators
- leaves are numbers

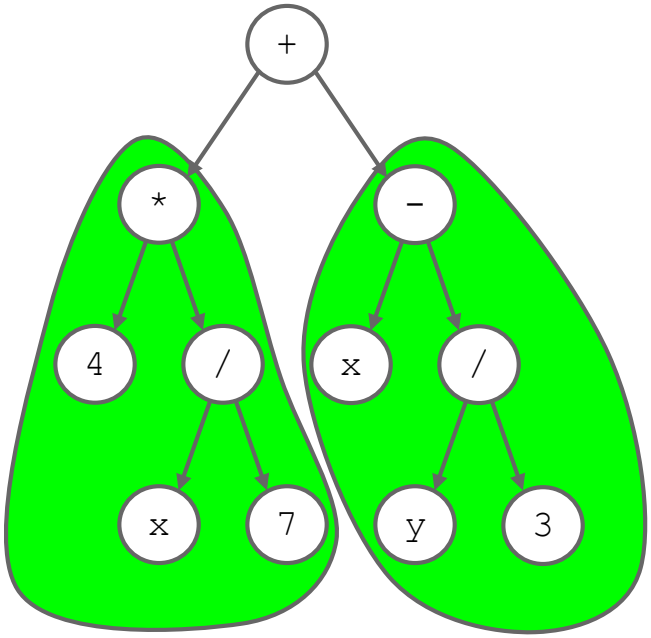
Thus, postfix expressions can be defined recursively, too.

E is a *postfix expression* when:

- *E* is a number, OR . . .
- *E* is two postfix expressions followed by a binary operator (+, −, *, /)

Algorithm to evaluate expression tree:

- bottom up



left exp'n tree right exp'n tree +

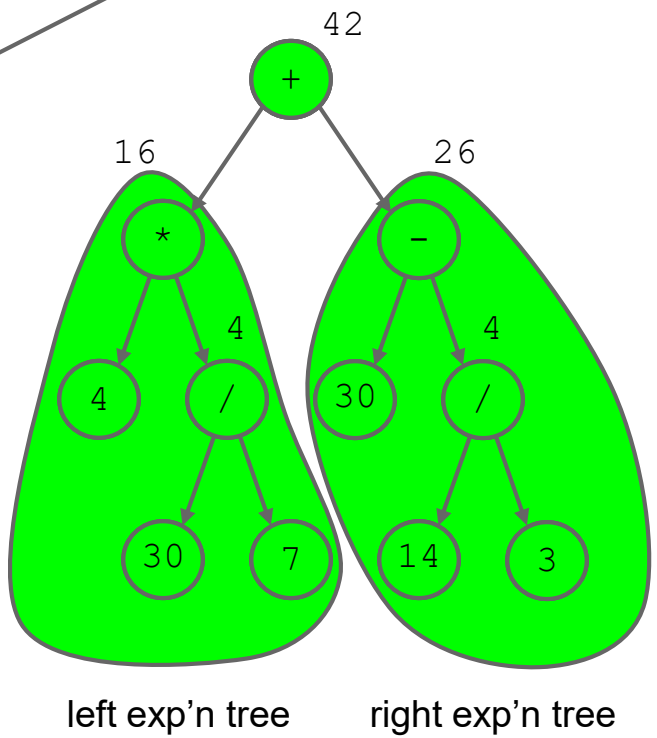
4 x 7 / * x y 3 / − +

1st operand 2nd operand

Tree Evaluation: Traversals

Algorithm to evaluate a tree rooted at vertex *x*. *recursion*

- If *x* has a number, then return that number
- If *x* has an operator, then:
  - evaluate the left subtree
  - evaluate the right subtree
  - return (left *op* right)



left exp'n tree right exp'n tree

Known as a *post-order traversal*

- evaluate the children first, then yourself
- follows the order: left → right → self

Other common traversals:

- self → left → right: *pre-order*
- left → self → right: *in-order*

Q. What's the in-order traversal?

Evaluated: `4 30 7 / * 30 14 3 / - +`

Q. What is this sequence?

Stack-Based Postfix Calculator

Use a Stack ADT to evaluate postfix.

Algorithm:

```
Create an empty stack S
while there is still input {
        if next input token is a number
                push the number to S
        if next input token is an operator {
                pop from S → b
                pop from S → a
                push (a op b) to S
        }
}
pop from S → result
```

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If any pop fails, then it's invalid postfix.

If S ends nonempty then it's invalid postfix.

Building Expression Trees from Postfix

Adapt postfix calculator algorithm to build trees from postfix.

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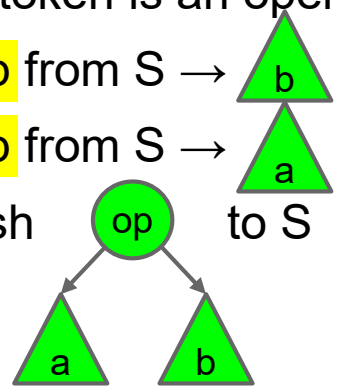

Example:

9 6 5 + 6 9 – * –

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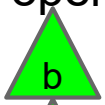
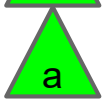
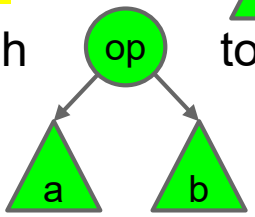
pop from S → 

pop from S → 

push 

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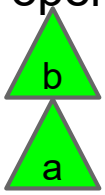
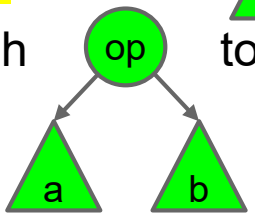
pop from S → 

push 

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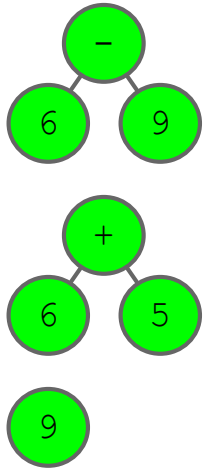

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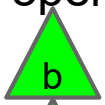
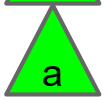
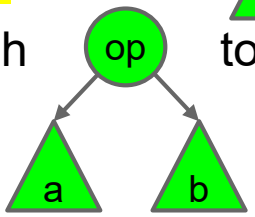
pop from S → 

pop from S → 

push  to S

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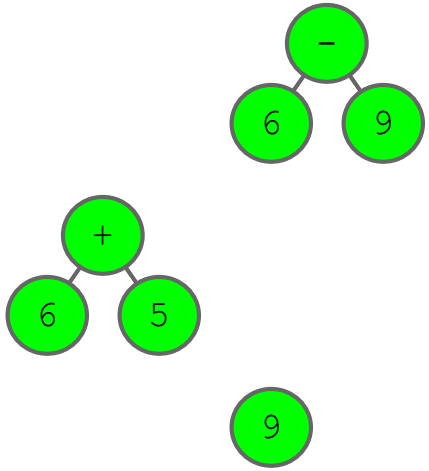

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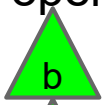
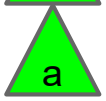
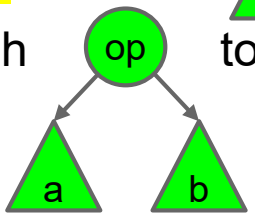
pop from S → 

pop from S → 

push 

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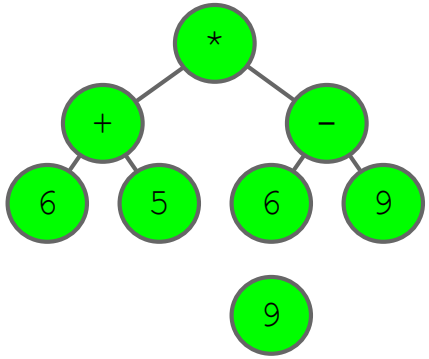

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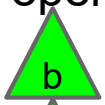
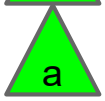
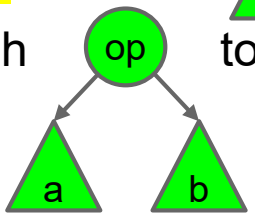
pop from S → 

pop from S → 

push  to S

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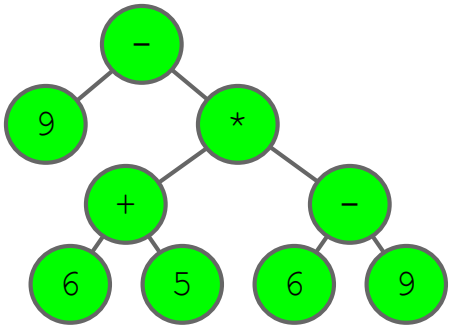

Expression Grammars

You can express the recursion using a *grammar*

Grammar for *E*:

- *E* → number
- *E* → *E E* operator

Production rules should be able to *derive* any valid expression, by stepwise substitution until no *E*'s remain.

Q. What's the grammar for infix?

E is a *postfix expression* when:

- *E* is a number, OR . . .
- *E* is two postfix expressions followed by a binary operator

E.g. Derive: `7 4 - 8 * 2 9 * +`

E ⇒ *E E* +
⇒ *E E* * *E* +
⇒ *E E* - *E* * *E* +
⇒ 7 *E* - *E* * *E* +
⇒ 7 4 - *E* * *E* +
⇒ 7 4 - 8 * *E* +
⇒ 7 4 - 8 * *E E* * +
⇒ 7 4 - 8 * 2 *E* * +
⇒ 7 4 - 8 * 2 9 * +